

05 - EBNF (Extended BNF) and Syntax Diagrams

EBNF is the same as BNF, with three additional meta-symbols:

- $\{ \}$ which indicates *0 or more*
- $[]$ which indicates *optional*
- $(\dots | \dots | \dots)$ which indicates sub-alternatives

EBNF has exactly the same expressive power as BNF.

But it is more convenient for many applications.

Converting from BNF to EBNF must be done precisely:

| BNF | EBNF |
|--|--|
| $\langle N \rangle ::= A AB$ | $\langle N \rangle ::= A [B]$ |
| $\langle Q \rangle ::= -\langle \text{num} \rangle \langle \text{num} \rangle$ | $\langle Q \rangle ::= [-] \langle \text{num} \rangle$ |
| $\langle P \rangle ::= \langle P \rangle A A$ | $\langle P \rangle ::= A \{ A \}$ |
| $\langle X \rangle ::= \langle X \rangle A \epsilon$ | $\langle X \rangle ::= \{ A \}$ |
| $\langle \text{blk} \rangle ::= \text{begin } \langle \text{sts} \rangle \text{ end}$ $\langle \text{sts} \rangle ::= \langle \text{cmd} \rangle \langle \text{cmd} \rangle ; \langle \text{sts} \rangle$ | $\langle \text{blk} \rangle ::= \text{begin } \langle \text{cmd} \rangle \{ ; \langle \text{cmd} \rangle \} \text{ end}$ |
| $\langle \text{nws} \rangle ::= +\langle \text{num} \rangle -\langle \text{num} \rangle$ | $\langle \text{nws} \rangle ::= (+ -) \langle \text{num} \rangle$ |
| $\langle \text{SN} \rangle ::= +\langle \text{num} \rangle -\langle \text{num} \rangle \langle \text{num} \rangle$ | $\langle \text{SN} \rangle ::= [(+ -)] \langle \text{num} \rangle$ |

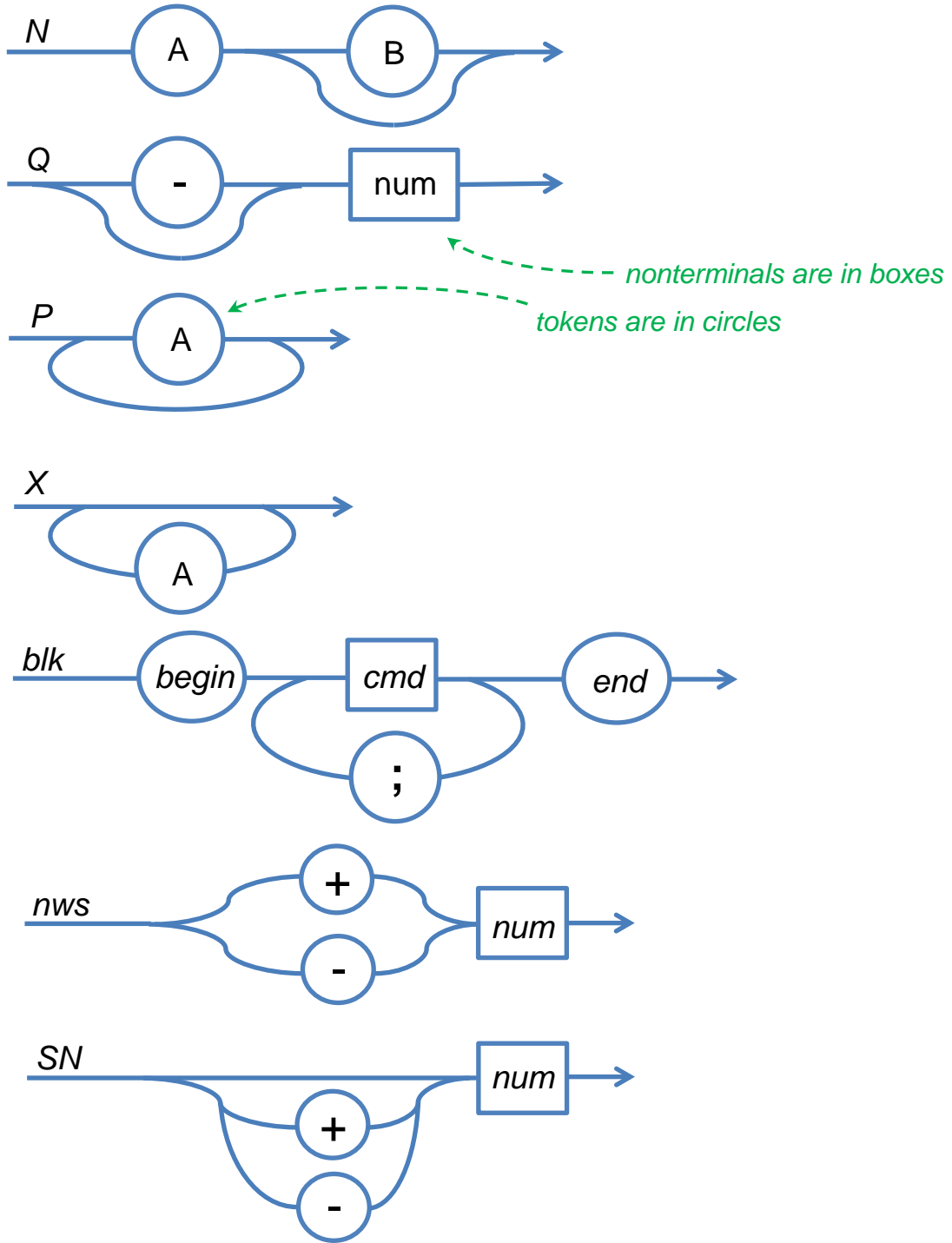
Some things to notice about the conversions to EBNF:

- most “or”s ($|$) have been removed, reducing the number of rules,
- redundant items are removed when all they do is specify options,
- most recursion has been removed and replaced with $\{ \}$ loops, and
- occurrences of the null string (ϵ) have been removed.

Conversion to EBNF makes it easier to draw Syntax Diagrams.

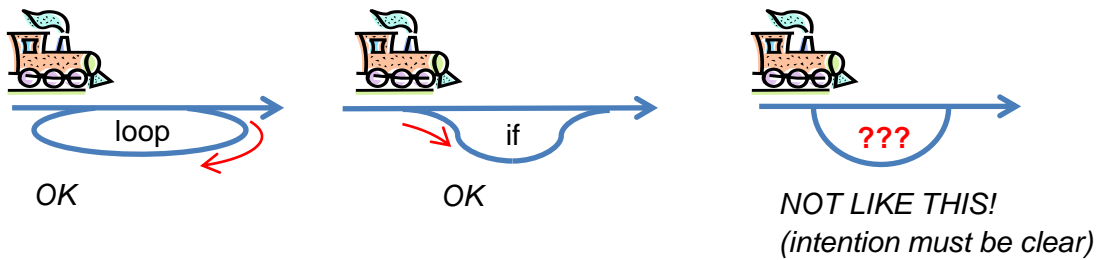
Later, we will use the Syntax Diagrams to write a recursive-descent parser.

Syntax Diagrams, sometimes called “Railroad Tracks”, are graphical representations of EBNF production rules. Here are syntax diagrams for each of the examples on the previous page:

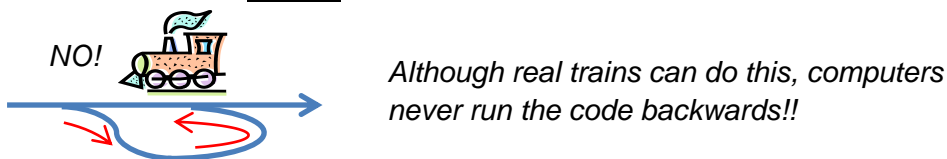


It can be helpful to imagine train tracks, to help in drawing them correctly:

- Control structures (curves and switches) should be very clear:



- The train must never “reverse directions”:



There are some common structures in programming languages.

Here is the correct way to draw them in BNF, EBNF, and Syntax Diagrams:

| | BNF | EBNF | Syntax Diagram |
|--|---|-------------------------|----------------|
| <i>A is optional</i> | $M ::= xxAxx \mid xxx$ | $M ::= xx[A]xx$ | |
| <i>A is required</i> | $M ::= xxAxx$ | $M ::= xxAxx$ | |
| <i>1 or more of A</i> | $M ::= MA \mid A$ | $M ::= A \{ A \}$ | |
| <i>0 or more of A</i> | $M ::= MA \mid \epsilon$ | $M ::= \{ A \}$ | |
| <i>1 or more of A with separators</i> | $M ::= M ; A \mid A$ | $M ::= A \{ ; A \}$ | |
| <i>1 or more of A with terminators</i> | $M ::= MA ; \mid A ;$ | $M ::= A \{ ; A ; \}$ | |
| <i>0 or more of A with separators</i> | $M ::= H \mid \epsilon$ $H ::= H ; A \mid A$ | $M ::= [A \{ ; A \}]$ | |
| <i>0 or more of A with terminators</i> | $M ::= MA ; \mid \epsilon$ | $M ::= \{ A ; \}$ | |